

LECTURE 21

Θ (Big Theta)

$\Theta(g(n))$ is the set of functions w/ the same order of growth as $g(n)$.

So $g(n)$ is asymptotically tight bound for $f(n)$

$\Theta(g(n)) = \{ f(n) \text{ there exist positive constants } c_1, c_2 \text{ and } n \text{ such that for all } n \geq n_0$
 $0 \leq c_1 g(n) \leq f(n) \leq c_2 g(n) \}$

$\rightarrow a(n)$

Ex. prove that $\underbrace{n^2 + 5n + 7}_{f(n)} = \Theta(n^2)$

Pf. when $n \geq 1$, then

$$1n^2 \leq \underbrace{n^2 + 5n + 7}_f \leq n^2 + 5n^2 + 7n^2 = 13n^2$$

g $C_2 g$

So, $n^2 + 5n + 7 = \Theta(n^2)$ by def. of Θ
with $n_0 = 1$, $C_1 = 1$ & $C_2 = 13$

Ex. prove that $n^5 + n^3 + 7n + 1 = \Theta(n^5)$

Properties of the Notations

Thm $f(n) = \Theta(g(n))$ iff $f(n) = O(g(n))$ &
 $f(n) = \Omega(g(n))$

Pf (Exc) Proof follows from the defns

Ex. Show that $\frac{1}{2}n^2 + 3n = \Theta(n^2)$

w/out using defn. of Θ

PF. IF $n \geq 1$, $\frac{1}{2}n^2 + 3n \leq \frac{1}{2}n^2 + 3n^2 = \frac{7}{2}n^2$

So, $\frac{1}{2}n^2 + 3n = O(n^2)$ — (1)

When $n \geq 0$, $\frac{1}{2}n^2 \leq \frac{1}{2}n^2 + 3n$

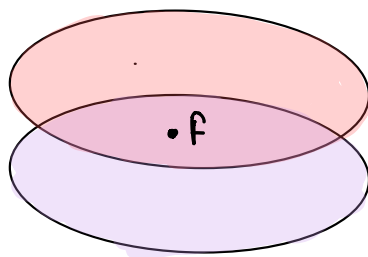
So, $\frac{1}{2}n^2 + 3n = \Omega(n^2)$


$\boxed{(1) + (2) + TH} \Rightarrow \frac{1}{2}n^2 + 3n = \Theta(n^2)$

Relations b/w

O , Θ

and Ω



 ; $\Omega(f)$

 : $\Theta(f)$

 : $O(f)$

Thm (Transitivity) IF $f(n) = O(g(n))$ and $g(n) = O(h(n))$, then $f(n) = O(h(n))$

Same for Big Θ & Ω



Pf. Exc.

Ex. Since $4n^2 + 3n + 17 \stackrel{\text{Exc.}}{\underset{\text{Def.}}{=}} O(n^3)$

and $n^3 \stackrel{\text{Exc.}}{\underset{\text{Def.}}{=}} O(n^4)$

We have $4n^2 + 3n + 17 \stackrel{\text{TH}}{=} O(n^4)$

Thm (Scaling by a constant)

If $f(n) = O(g(n))$, then $Kf(n) = O(g(n))$
for any $K > 0$

Pf. Assume that $f(n) = O(g(n))$
then by def. there exist positive
constants c and n_0 such that

$f(n) \leq cg(n)$ for all $n \geq n_0$

let $K > 0$, then $Kf(n) \leq \underbrace{(Kc)}_{c'} g(n)$
for all $n \geq 0$

Thus $Kf(n) = O(g(n))$

FACT ^{from Calculus} - let a, b and x be positive real #s
with $a, b \neq 1$ Then $\log_a^x = \frac{\log_b^x}{\log_b^a}$

Ex let $a, b \neq 1$ be positive #s then
 $\log_a^n = O(\log_b^n)$

Thm (Sums) If $f_1(n) = O(g_1(n))$ and $f_2(n) = O(g_2(n))$
then $f_1(n) + f_2(n) = O(g_1(n) + g_2(n))$
 $= O(\max\{g_1(n), g_2(n)\})$
Same for Θ & Ω

Pf. We prove that the result for O
notation

$$f_i(n) = O(g_i(n)) \text{ for } i=1,2$$

Def - There exist + constants c_i and n_i ;
such that $f_i(n) \leq c_i g_i(n)$, for all
 $n \geq n_i$ ($i=1,2$)

$$\text{let } n_0 = \max\{n_1, n_2\}$$

If $n \geq n_0$, we have

$$\begin{aligned} f_1(n) + f_2(n) &\leq c_1 g_1(n) + c_2 g_2(n) \\ &\leq \max\{c_1, c_2\} g_1(n) \\ &\quad + \max\{c_1, c_2\} g_2(n) \end{aligned}$$

$$= \max\{C_1, C_2\} (g_1(n) + g_2(n))$$

constant

$$\leq \max\{C_1, C_2\} (2 * \max\{g_1(n), g_2(n)\})$$

constant

Thus, $f_1 + f_2 \in O(g_1 + g_2)$

also $f_1 + f_2 \in O(\max_i g_i)$

Ex. previously we proved
 $n^2 + n = O(n^2)$ and $3n^3 \dots = O(n^3)$

Then
 $(n^2 + n) + (3n^3 - 2n^2 + 13n - 15) = O(n^3)$
max $\{n^2, n^3\}$

Thm (Products)

if $f_1(n) = O(g_1(n))$ and $f_2(n) = O(g_2(n))$
 then $f_1(n) f_2(n) = O(g_1(n) g_2(n))$

And for Θ and \sim

$$(n^2 + n)(3n^3 - 2n^2 + 13n - 15) = O(n^5)$$

Thm (Symmetry) $f(n) = O(g(n))$
IFF $g(n) = \Omega(f(n))$
Thm (Reflexive) $f(n) = O(f(n))$

Homework:

<https://u.osu.edu/alzalg.1/files/2019/10/HW9W.pdf>