

Lecture 20

Running Time (Review of Lecture 19 material)

Ex. Find UB & LB of the running time of the alg. represented by

$$\sum_{i=1}^n \sum_{j=1}^{i^2} \sum_{k=j}^{i^2} c = \sum_{i=1}^n \sum_{j=1}^{i^2} c(i^2 - j)$$

UB:
$$\sum_{i=1}^n \sum_{j=1}^{i^2} c(i^2 - j) \leq \sum_{i=1}^n \sum_{j=1}^{i^2} c(i^2)$$

$$= \sum_{i=1}^n c(i^2)(i^2)$$

$$= \sum_{i=1}^n ci^4 \leq \sum_{i=1}^n cn^4 = c(n^4)(n)$$

$$= cn^5$$

LB:
$$\sum_{i=1}^n \sum_{j=1}^{i^2} c(i^2 - j) \geq \sum_{i=1}^n \sum_{j=1}^{i^2/2} c(i^2 - j)$$

① *Split the sum*

$$\geq \sum_{i=1}^n \sum_{j=1}^{i^2/2} c(i^2 - i^2/2)$$

$$= \sum_{i=1}^n c(i^2/2)(i^2/2)$$

$$\begin{aligned}
 & i=1 \quad \dots \quad \dots \quad \dots \\
 & = \sum_{i=1}^n c i^4 / 4 \geq \sum_{i=n/2}^n c i^4 / 4 \geq \sum_{i=n/2}^n \frac{c}{4} \left(\frac{n}{2}\right)^4 \\
 & = \frac{c}{64} \sum_{i=n/2}^n n^4 = \frac{c}{64} n^4 \left(\frac{n}{2}\right) \\
 & = \frac{c n^5}{128}
 \end{aligned}$$

you get something smaller

The running time is

$$c' n^5 \text{ where } \frac{c}{128} \leq c' \leq c$$

Asymptotic Notation

- behavior of a function as the input approaches ∞

NOTATION

- Big O - let f be a nonnegative function

$f(n)$ is O of $g(n)$, written as $f(n) = O(g(n))$, iff there exist positive constants c and n_0 such that

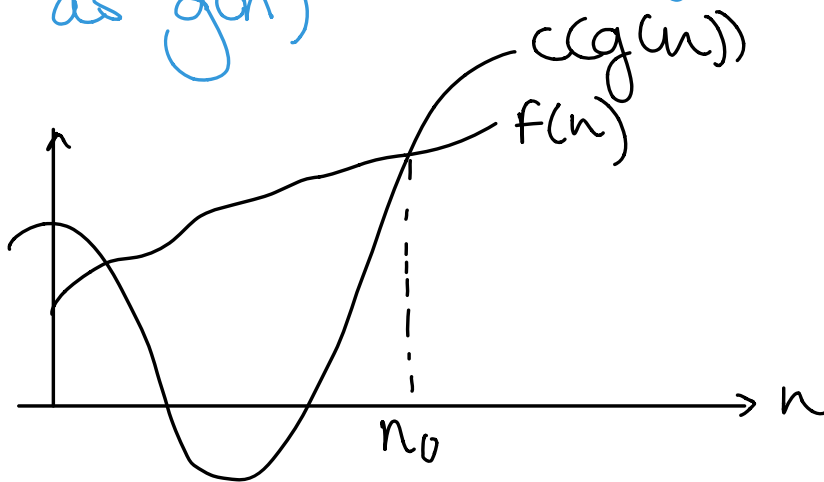
$$f(n) \leq c g(n) \text{ for all } n \geq n_0.$$

If $f(n) = O(g(n))$, the $f(n)$ grows no faster than $g(n)$

* $g(n)$ is an asymptotic upper bound (or just UB) on $f(n)$

$$O(g(n)) = \{f(n) : \text{(there exist positive constants } c \text{ and } n \text{ such that } \underline{0 \leq f(n) \leq c g(n)} \text{ for all } \underline{n \geq n_0})}\}$$

The set of functions with smaller or same order of growth as $g(n)$



$$f(n) \in O(g(n))$$

$$\left. \begin{array}{l} A \leq B \\ C \leq D \end{array} \right\} \Rightarrow A+C \leq B+D$$

Ex. prove that $\underbrace{(n^2+n)}_{f(n)} = O(\underbrace{n^3}_{g(n)})$

Sol: For $n \geq \underline{1}$, we have

$$n^2 \leq n^3 \text{ and } n \leq n^3$$

Then $\underbrace{(n^2+n)}_{f(n)} \leq \underbrace{n^3}_{\downarrow} + \underbrace{n^3}_{\downarrow} = 2n^3$
raise the powers

Ex: prove that $3n^3 - 2n^2 + 13n - 15 = O(n^3)$

Pf: For $n \geq 1$, we have → remove negatives

$$\begin{aligned} 3n^3 - 2n^2 + 13n - 15 &\leq 3n^3 + 13n \\ &\leq 3n^3 + 13n^3 \leq 16n^3 \end{aligned}$$

Ex. Prove that $n^{7/2} + n^3 \log n = O(n^4)$

Pf: If $n \geq 1$, we have

$$n^{7/2} \leq n^4$$

$$\underline{\log n \leq n}$$

$$\log n \leq n \Rightarrow n^3 \log n \leq n^4$$

$$n^{7/2} + n^3 \log n \leq n^4 + n^4 = 2n^4$$